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# DEPARTMENT OF COMPUTER 

 SCIENCE AND ARTIFICIAL INTELLIGENCE (CS\&AI)COMPUTER
ORGANIZATION AND MICROPROCESSOR LAB MANUAL (PC451AD) B.E IV SEM (2021-2022)

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## OVERVIEW OF 8086 MICROPROCESSOR ARCHITECTURE OF 8086:

As shown in the below figure, the 8086 CPU is divided into two independent functional parts. Dividing the work between these two units' speeds up processing.

- Bus Interface Unit(BIU)
- Execution Unit(EU)



## The Execution Unit (EU):

- The execution unit of the 8086 tells the BIU where to fetch instructions or data from, decodes instructions, and executes instructions.
- The EU contains control circuitry, which directs internal operations. A decoder in the EU translates instructions fetched from memory into a series of actions, which the EU carries out.
- The EU has a 16-bit arithmetic logic unit (ALU) which can add, subtract, AND, OR, OR, increment, decrement, complement or shift binary numbers.
- The main functions of EU are:
o Decoding of Instructions
o Execution of instructions


## Steps:

- EU extracts instructions from top of queue in BIU
- Decode the instructions
- Generates operands if necessary
- Passes operands to BIU \& requests it to perform read or write bus cycles to memory or I/O
- Perform the operation specified by the instruction on operands


## Bus Interface Unit (BIU):

- The BIU sends out addresses, fetches instructions from memory, reads data from ports and memory, and writes data to ports and memory.
- In simple words, the BIU handles all transfers of data and addresses on the buses for the execution unit.


## 8086 has Pipelining Architecture:

- While the EU is decoding an instruction or executing an instruction, which does not require use of the buses, the BIU fetches up to six instruction bytes for the following instructions.
- The BIU stores these pre-fetched bytes in a first-in-first-out register set called a queue. When the EU is ready for its next instruction from the queue in the BIU. This is much faster than sending out an address to the system memory and waiting for memory to send back the next instruction byte or bytes.
- Except in the case of JMP and CALL instructions, where the queue must be dumped and then reloaded starting from a new address, this pre-fetch and queue scheme greatly speeds up processing.
- Fetching the next instruction while the current instruction executes is called pipelining.


## REGISTER ORGANIZATION:

- 8086 has a powerful set of registers known as general purpose registers and special purpose registers. All of them are 16 -bit registers.
- The 8086 registers are classified into the following types:
o General Data Registers
o Segment Registers
o Pointers and Index Registers
o Flag Register
The register set of 8086 can be categorized into 4 different groups. The register organization of 8086 is shown in the figure.

| $A X$ | $A H$ | $A L$ |
| :--- | :--- | :---: |
| $B X$ | $B H$ | $B L$ |
| $C X$ | $C H$ | $C L$ |
| $D X$ | $D H$ | $D L$ |
|  |  |  |

General data registers


Segment Registers


Pointers and index registers

Register organization of $\$ 086$

## General Data Registers:

- The registers $A X, B X, C X$ and $D X$ are the general purpose 16-bit registers.
- $A X$ is used as 16 -bit accumulator. The lower 8 -bit is designated as AL and higher 8 -bit is designated as AH . AL can be used as an 8-bit accumulator for 8-bit operation.
- All data register can be used as either 16 bit or 8 bit. BX is a 16 bit register, but BL indicates the lower 8-bit of $B X$ and $B H$ indicates the higher 8 -bit of $B X$.
- The register $B X$ is used as offset storage for forming physical address in case of certain addressing modes.
- The register CX is used default counter in case of string and loop instructions
- DX register is a general purpose register which may be used as an implicit operand or destination in case of a few instructions.


## Segment Registers:

- The 8086 architecture uses the concept of segmented memory. 8086 able to address a memory capacity of 1 megabyte and it is byte organized. This 1 megabyte memory is divided into 16 logical segments. Each segment contains 64 kbytes of memory. There are 4 segment registers.
o Code Segment Register(CS):is used for addressing memory location in the code segment of the memory, where the executable program is stored.
o Data Segment Register(DS): points to the data segment of the memory where the data is stored.
o Extra Segment Register(ES):also refers to a segment in the memory which is another data segment in the memory.
o Stack Segment Register(SS):is used for addressing stack segment of the memory. The stack segment is that segment of memory which is used to store stack data.
- While addressing any location in the memory bank, the physical address is calculated from two parts:

Physical address= segment address + offset address

- The first is segment address, the segment registers contain 16-bit segment base addresses, related to different segments. The second part is the offset value in that segment.


## Pointers and Index Registers:

- The pointers registers contain offset within the particular segments. The index and pointer registers are given below:
- IP(Instruction pointer)-store memory location of next instruction to be executed. The pointer register IP contains offset within the code segment.
- BP(Base pointer)-The pointer register BP contains offset within the data segment.
- SP(Stack pointer)- The pointer register SP contains offset within the stack segment.
- The index registers are used as general purpose registers as well as for offset storage in case of indexed, base indexed and relative base indexed addressing modes
- SI(Source index)- The register SI is used to store the offset of source data in data segment.
- Dl(Destination index)- The register DI is used to store the offset of destination in data or extra segment.


## 8086 flag register and its functions:

- The 8086 flag register contents indicate the results of computation in the ALU. It also contains some flag bits to control the CPU operations.
- A 16 bit flag register is used in 8086 . It is divided into two parts .
- Condition code or status flags
- Machine control flags
- The condition code flag register is the lower byte of the 16 -bitflag register. The condition code flag register is identical to 8085 flag register, with an additional overflow flag
- The control flag register is the higher byte of the flag register. It contains three flags namely direction flag (D), interrupt flag (I) and trap flag (T).

- SF- Sign Flag: This flag is set, when the result of any computation is negative. For signed computations the sign flag equals the MSB of the result.
- ZF-Zero Flag: This flag is set, if the result of the computation or comparison performed by the previous instruction is zero.
- PF-Parity Flag: This flag is set to 1 , if the lower byte of the result contains even number of 1 's
- CF-Carry Flag: This flag is set, when there is a carry out of MSB in case of addition or a borrow in case of subtraction.
- AF(Auxiliary Carry Flag): This is set, if there is a carry from the lowest nibble, i.e, bit three during addition, or borrow for the lowest nibble, i.e, bit three, during subtraction.
- OF- Over flow Flag: This flag is set, if an overflow occurs, i.e, if the result of a signed operation is large enough to accommodate in a destination register. The result is of more than 7 -bits in size in case of 8 -bit signed operation and more than 15 -bits in size in case of 16 -bit sign operations, and then the overflow will be set.
- TF- Tarp Flag: If this flag is set, the processor enters the single step execution mode. The processor executes the current instruction and the control is transferred to the Trap interrupt service routine.
- IF- Interrupt Flag: If this flag is set, the mask able interrupts are recognized by the CPU, otherwise they are ignored.
- D- Direction Flag: This is used by string manipulation instructions. If this flag bit is ' 0 ', the string is processed beginning from the lowest address to the highest address, i.e., auto incrementing mode. Otherwise, the string is processed from the highest address towards the lowest address, i.e., auto decrementing mode.


## Addressing modes of 8086:

- Addressing mode indicates a way of locating data or operands.
- The addressing modes describe the types of operands and the way they are accessed for executing an instruction.
- According to the flow of instruction execution, the instructions may be categorized as Sequential control flow instructions and Control transfer instructions
- Sequential control flow instructions are the instructions, which after execution, transfer control to the next instruction appearing immediately after it (in the sequence) in the program. For example, the arithmetic, logic, data transfer and processor control instructions are sequential control flow instructions.
- The control transfer instructions, on the other hand, transfer control to some predefined address or the address somehow specified in the instruction, after their execution. For example, INT, CALL, RET and JUMP instructions fall under this category.
- The addressing modes for sequential control transfer instructions are:

1. Immediate: In this type of addressing, immediate data is a part of instruction and appears in the form of successive byte or bytes. Ex: MOV AX,0005H
In the above example, 0005 H is the immediate data. The immediate data may be 8 -bit or 16 -bit in size.
2. Direct: In the direct addressing mode a 16-bit memory address (offset) is directly specified in the instruction as a part of it.
Ex: MOV AX, $[5000 \mathrm{H}]$
Here, data resides in a memory location in the data segment, whose effective address may be completed using 5000 H as the offset address and content of DS as segment address. The effective address here, is 10 H * DS +5000 H .
3. Register: In register addressing mode, the data is stored in a register and is referred using the particular register. All the registers, except IP, may be used in this mode.
Ex: MOV BX, AX
4. Register Indirect: Sometimes, the address of the memory location, which contains data or operand, is determined in an indirect way, using the offset register. This mode of addressing is known as register indirect mode. In this addressing mode, the offset address of data is in either BX or SI or DI register. The default segment is either DS or ES. The data is supposed to be available at the address pointed to by the content of any of the above registers in the default data segment.
Ex: MOV AX, $[B X]$
Here, data is present in a memory location in DS whose offset address is in BX. The effective address of the data is given as $10 \mathrm{H}^{*} \mathrm{DS}+[\mathrm{BX}]$.
5. Indexed: In this addressing mode, offset of the operand is stored in one of the index registers. DS and ES are the default segments for index registers, SI and DI respectively. This is a special case of register indirect addressing mode.
Ex: MOV AX, [SI]
Here, data is available at an offset address stored in SI in DS. The effective address, in this case, is computed as $10 * \mathrm{DS}+[\mathrm{SI}]$.
6. Register Relative: In this addressing mode, the data is available at an effective address formed by adding an 8-bit or 16-bit displacement with the content of anyone of the registers BX , $B P, S I$ and $D I$ in the default (either DS or ES) segment. Ex: MOV AX, $50 \mathrm{H}[B X]$ Here, the effective address is given as $10 \mathrm{H} * \mathrm{DS}+50 \mathrm{H}+[\mathrm{BX}]$
7. Based Indexed: The effective address of data is formed, in this addressing mode, by adding content of a base register (any one of BX or BP) to the content of an index register (any one of SI or DI). The default segment register may be ES or DS.
Ex: MOV AX, [BX][SI]
Here, BX is the base register and S is the index register the effective address is computed as $10 H^{*} D S+[B X]+[S t]$.
8. Relative Based Indexed: The effective address is formed by adding an 8 or 16 -bit displacement with the sum of the contents of any one of the base register (BX or BP) and any one of the index register, in a default segment.
Ex: MOV AX, $50 \mathrm{H}[\mathrm{BX}][\mathrm{SI}]$
Here, 50 H is an immediate displacement, BX is base register and SI is an index register the effective address of data is computed as
10 H * DS + [BX] + [SI] + 50 H
.9. Intrasegment Direct Mode: In this mode, the address to which the control is to be transferred lies in the same segment in which the control transfer instruction lies and appears directly in the instruction as an immediate displacement value. In this addressing mode, the displacement is computed relative to the content of the instruction pointer IP.

The effective address to which the control will be transferred is given by the sum of 8 or 16 -bit displacement and current content of IP. In the case of jump instruction, if the signed displacement (d) is of 8 -bits (i.e $-128<\mathrm{d}<+128$ ) we term it as short jump and if it is of 16 -bits (i.e-32, $768<\mathrm{d}<+32,768$ ) it is termed as long jump.
10. Intrasegment Indirect Mode: In this mode, the displacement to which the control is to be transferred, is in the same segment in which the control transfer instruction lies, but it is passed to the instruction indirectly. Here, the branch address is found as the content of a register or a memory location. This addressing mode may be used in unconditional branch instructions.
11. Intersegment Direct: In this mode, the address to which the control is to be transferred is in a different segment. This addressing mode provides a means of branching from one code segment to another code segment. Here, the CS and IP of the destination address are specified directly in the instruction.
12. Intersegment Indirect: In this mode, the address to which the control is to be transferred lies in a different segment and it is passed to the instruction indirectly, i.e contents of a memory block containing four bytes, i.e IP (LSB), IP(MSB), CS(LSB) and CS (MSB) sequentially. The starting address of the memory block may be referred using any of the addressing modes, except immediate mode.

## Pin Diagram of 8086:

- The 8086 is a 16 -bit microprocessor. This microprocessor operates in single processor or multiprocessor configurations to achieve high performance.
- The pin configuration of 8086 is shown in the figure. Some of the pins serve a particular function in minimum mode (single processor mode) and others function in maximum mode (multiprocessor mode).
- The 8086 signals are categorized into 3 types:
1.Common signals for both minimum mode and maximum mode.

2. Special signals which are meant only for minimum mode
3. Special signals which are meant only for maximum mode

- Common Signals for both Minimum mode and Maximum mode:

1. AD7-AD0: The address/ data bus lines are the multiplexed address data bus and contain the right most eight bit of memory address or data. The address and data bits are separated by using ALE signal.
2. AD15- AD8: The address/data bus lines compose the upper multiplexed address/data bus. This lines contain address bit A15 A8 or data bus D15 D8. The address and data bits are separated by using ALE signal.
3. A19 / S6- A18 / S3: The address/status bus bits are multiplexed to provide address signals A19 /A16 and also status bits S6 /S3. The address bits are separated from the status bits using the ALE signals. The status bit S6 is always a logic 0 , bit S 5 indicates the condition of the interrupt flag bit. The S4 and S3 being used for memory access indicate which segment register is presently.

| S4 | S3 | Type of segment register used |
| :--- | :--- | :--- |
| 0 | 0 | Extra segment |
| 0 | 1 | Stack segment |
| 1 | 0 | Code or no segment |
| 1 | 1 | Data Segment |

MAX MODE


HOLD
HLDA
$\overline{\text { WR }}$
M/历
DT/R
DEN
ALE
INTA
4. BHE / S7:The bus high enable (BHE) signal is used to indicate the transfer of data over the higher order (D15-D8) data bus. It goes low for the data transfer over D15-D8 and is used to derive chip select of odd address memory bank or peripherals.

5. $\boldsymbol{\operatorname { R D } ( \text { Read ): whenever the read signal is at logic } 0 \text { , the data bus receiyés the data from the }}$ memory or I/O devices connected to the system
6. READY: This is the acknowledgement from the slow devices or memory that they have completed the data transfer operation. This signal is active high.
7. INTR(Interrupt Request): Interrupt request is used to request a hardware interrupt of INTR is held high when interrupt enable flag is set, the 8086 enters an interrupt acknowledgement cycle after the current instruction has completed its execution.
8. TEST : This input is tested by "WAIT" instruction. If the TEST input goes low; execution will continue. Else the processor remains in an idle state.
9. NMI(Non-maskable Interrupt): The non-maskable interrupt input is similar to INTR except that the NMI interrupt does not check for interrupt enable flag is at logic 1 , i.e, NMI is not maskable internally by software. If NML is activated, the interrupt input uses interrupt vector 2.
10. RESET: The reset input causes the microprocessor to reset itself. When 8086 reset, it restarts the execution from memory location FFFFOH. The reset signal is active high and must be active for at least four clock cycles.
11. CLK: Clock input: The clock input signal provides the basic timing input signal for processor and bus control operation. It is asymmetric square wave with $33 \%$ duty cycle.
12. VCC $(+5 \mathrm{~V})$ : Power supply for the operation of the internal circuit
13. GND: Ground for the internal circuit
14. MN / MX : The minimum/maximum mode signal to select the mode of operation either in minimum or maximum mode configuration. Logic 1 indicates minimum mode.

## Minimum mode Signals:

The following signals are for minimum mode operation of 8086 .

1. $\mathbf{M} / \mathrm{IO}$ ( Memory/IO):M / IO signal selects either memory operation or I/O operation. This line indicates that the microprocessor address bus contains either a memory address or an I/O port address. Signal high at this pin indicates a memory operation. This line is logically equivalent to S 2 in maximum mode.
2. INTA(Interrupt acknowledge): The interrupt acknowledge signal is a response to the INTR input signal. The INTA signal is normally used to gate the interrupt vector number onto the data bus in response to an interrupt request.
3. ALE(Address Latch Enable): This output signal indicates the availability of valid address on the address/data bus, and is connected to latch enable input of latches.
4. DT / R ( Data transmit/Receive): This output signal is used to decide the direction of date flow through the bi-directional buffer. DT / R 1 Indicates transmitting and DT/R 0 indicates receiving the data.
5. DEN(Data Enable): Data bus enable signal indicates the availability of valid data over the address/data lines.
6. (Write): whenever the write signal is at logic 0 , the data bus transmits the data to the memory or I/O devices connected to the system.
7. HOLD: The hold input request a direct memory access (DMA). If the hold signal is at logic 1 , the micro process stops its normal execution and places its address, data and control bus at the high impedance state.
8. HLDA: Hold acknowledgement indicates that 8086 has entered into the hold state.

## Maximum mode signal:

- The following signals are for maximum mode operation of 8086.

1. $\underline{\mathbf{S 2}, \mathbf{S 1}, \mathbf{S} 0(S t a t u s}$ lines): These are the status lines that reflect the type of operation being carried out by the processor. These status lines are encoded as follow

| S2 | S1 | S0 | Function |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | Interrupt Acknowledge |
| 0 | 0 | 1 | Read I/O port |
| 0 | 1 | 0 | Write I/O port |
| 0 | 1 | 1 | Halt |
| 1 | 0 | 0 | Code Access |
| 1 | 0 | 1 | Read Memory |
| 1 | 1 | 0 | Write Memory |
| 1 | 1 | 1 | Passive (In active) |

2. LOCK : The lock output is used to lock peripherals off the system, i.e, the other system bus masters will be prevented from gaining the system bus.
3. QS1 and QS0 (Queue status): The queue status bits shows the status of the internal instruction queue. The encoding of these signals is as follows

| QS1 | QS0 | Function |
| :--- | :--- | :--- |
| 0 | 0 | No operation,queue is idle |
| 0 | 1 | First byte of opcode |
| 1 | 0 | Queue is empty |
| 1 | 1 | Subsequent byte of opcode |

## 4. RQ / GT1 and RQ / GT 0 (request/Grant): The request/grant pins are used by other

 local bus masters to force the processor to release the local bus at the end of the processors current bus cycle. These lines are bi- directional and are used to both request and grant a DMA operation. RQ / GT 0 is having higher priority than RQ / GT1

## EXPERIMENT-1

AIM:- Write a Program to add two 8 bit numbers


## EXPERIMENT-2

AIM:- Write a Program to Subtract Two 8 bit numbers


## EXPERIMENT-3

AIM:- Write a Program to Multiply two 8 bit numbers


## EXPERIMENT-4

AIM:- Write a Program to Divide two 8 bit numbers


## EXPERIMENT-5

AIM:- Write a Program to add two 16 bit numbers


## EXPERIMENT-6

AIM:- Write a Program to Subtract Two 16 bit numbers


## EXPERIMENT-

AIM:- Write a Program to Multiply two 16 bit numbers


## EXPERIMENT-

AIM:- Write a Program to Divide two 16 bit numbers


## EXPERIMENT-

AIM:- Write a Program to add n 16 -Bit numbers.

RESULTS:-

| INPUT |  |  | OUTPUT |  |
| :---: | :---: | :---: | :---: | :---: |
| A050 |  | AX |  |  |
| A052 |  | DX |  |  |
| A054 |  | CX |  |  |
| A056 |  | SI |  |  |
| A058 |  |  |  |  |

## EXPERIMENT-24

AIM:- Write a Program to find the largest of two 8-bit numbers


## EXPERIMENT-25

AIM:- Write a Program to find the smallest of two 8-bit numbers PROGRAM:-

| 8000 | MOV | SI, A050 |
| :--- | :--- | :--- |
| 8003 | MOV | AL, [SI] |
| 8005 | INC | SI |
| 8006 | MOV | BL, [SI] |
| 8008 | CMP | AL,BL |
| 800A | JC | 800 E (Small) |
| 800C | MOV | AL,BL |
| Small: 800 E | INT | 03 |



## EXPERIMENT-26

AIM:- Write a Program to find the largest of two 16-BIT numbers.


## EXPERIMENT-27

AIM:- Write a Program to find the smallest of two 16-BIT numbers.


## EXPERIMENT-28

AIM:- Write a Program to find the largest of $n 8$-BIT numbers. PROGRAM:-

| 8000 | MOV | CL,04 |
| ---: | :--- | :--- |
| 8002 | MOV | SI, A050 |
| 8005 | MOV | AL, [SI] |
| Back: 8007 | INC | SI |
| 8008 | MOV | BL,[SI] |
| 800A | CMP | AL,BL |
| 800C | JNC | 8010 (Large) |
| Large: 80010 | MOV | AL,BL |
| 8012 | LOOP | 8007 (Back) |
|  |  | 03 |
|  | ORG | A050 |
|  | DB | -- |
|  | DB | -- |
|  | DB | -- |
|  | DB | -- |
| RESULTS:- | DB | - |


| INPUT |  | OUTPUT |  |
| :---: | :---: | :---: | :---: |
| A050 |  | AX |  |
| A051 |  | BX |  |
| A052 |  | CX |  |
| A053 |  | SI |  |
| A054 |  |  |  |

## EXPERIMENT-29

AIM:- Write a Program to find the smallest of n 8 -BIT numbers. PROGRAM:-

| 8000 | MOV | CL,04 |
| ---: | :--- | :--- |
| 8002 | MOV | SI, A050 |
| 8005 | MOV | AL, [SI] |
| Back: 8007 | INC | SI |
| 8008 | MOV | BL,[SI] |
| 800A | CMP | AL,BL |
| 800C | JC | 8010 (Small) |
| 800E | MOV | AL,BL |
| Small: 8010 | LOOP | 8007 (Back) |
| 8012 | INT | 03 |
|  |  |  |
|  | ORG | A050 |
|  | DB | -- |
|  | DB | -- |
|  | DB | -- |
|  | DB | -- |
| RESULTS:- | DB | -- |


| INPUT |  | OUTPUT |  |
| :---: | :---: | :---: | :---: |
| A050 |  | AX |  |
| A051 |  | BX |  |
| A052 |  | CX |  |
| A053 |  | SI |  |
| A054 |  |  |  |

## EXPERIMENT-30

AIM:- Write a Program to find the largest of n 16 -BIT numbers.

## PROGRAM:-



| INPUT |  | OUTPUT |  |
| :---: | :---: | :---: | :---: |
| A050 |  | AX |  |
| A052 |  | BX |  |
| A054 |  | CX |  |
| A056 |  | SI |  |
| A058 |  |  |  |

## EXPERIMENT-31

AIM:- Write a Program to find the smallest of n 16 -BIT numbers.
PROGRAM:-

| 8000 | MOV | CL,04 |
| ---: | :--- | :--- |
| 8002 | MOV | SI, A050 |
| 8005 | MOV | AX, [SI] |
| Back: 8007 | INC | SI |
| 8008 | INC | SI |
| 8009 | MOV | BX,[SI] |
| 800B | CMP | AX,BX |
| 800D | JC | 8011 (Small) |
| 800F | MOV | AX,BX |
| Small: 8011 | LOOP | 8007 (Back) |
| 8013 | INT | 03 |
|  |  |  |
|  | ORG | A050 |
|  | DW | ---- |
|  | DW | ---- |
|  | DW | ---- |
|  | DW | ---- |
|  | DW | ---- |


RESULTS:-

| INPUT |  | OUTPUT |  |
| :---: | :---: | :---: | :---: |
| A050 |  | AX |  |
| A052 |  | BX |  |
| A054 |  | CX |  |
| A056 |  | SI |  |
| A058 |  |  |  |

## EXPERIMENT-32

AIM:- Write a Program to perform Block Transfer PROGRAM:-


## EXPERIMENT-19

AIM:- Write a Program to sort n numbers

## PROGRAM:-

| 8000 | MOV | CH,04 |
| ---: | :--- | :--- |
| Loop1:8002 | MOV | CL,04 |
| 8004 | MOV | SI,A050 |
| Loop2:8007 | MOV | AL, [SI] |
| 8009 | INC | SI |
| 800A | CMP | AL,[SI] |
| 800C | JNC | 8014 (Next) |
| 800E | XCHG | AL,[SI] |
| 8010 | DEC | SI |
| 8011 | MOV | $[$ SI],AL |
| 8013 | INC | SI |
| Next: 8014 | DEC | CL |
| 8016 | JNZ | 8007 (Loop 2) |
| 8018 | DEC | CH |
| 801 A | JNZ | 8002 (Loop 1) |
| 801 C | INT | 03 |
|  |  |  |
|  | ORG | A050 |
|  | DB | -- |
|  | DB | -- |
|  | DB | -- |
|  | DB | -- |
|  | DB | -- |



## RESULTS:-

| INPUT |  | DATA | OUTPUT |  |
| :---: | :---: | :---: | :---: | :---: |
| A050 |  | A050 |  |  |
| A051 |  | A051 |  |  |
| A052 |  | A052 |  |  |
| A053 |  | A053 |  |  |
| A054 |  | A054 |  |  |

## EXPERIMENT-35

AIM:- Write a Program to search a 8 bit number PROGRAM:-


## EXPERIMENT-21

AIM:- Write a Program to find number of 1's in an 8 bit number PROGRAM:-


## EXPERIMENT-22

AIM:- Write a Program to find 2's complement of a number


## EXPERIMENT-23

AIM:- Write a Program to find negative of a number using NEG


